

# Deduction, Induction, Statistics, and Topology

{Kevin T. Kelly, Konstantin Genin}
Carnegie Mellon University

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# INDUCTIVE VS. DEDUCTIVE INFERENCE

- All the objects of human ... enquiry may naturally be divided into two kinds, to wit,
  - 1. Relations of Ideas, and
  - 2. Matters of Fact.

David Hume, Enquiry, Section IV, Part 1.

- Any ... inference in science belongs to one of two kinds:
  - either it yields certainty in the sense that the conclusion is necessarily true, provided that the premises are true,
  - 2. or it does not.
- The first kind is ... deductive inference ....
- The second kind will ... be called 'inductive inference'.
- R. Carnap, The Continuum of Inductive Methods, 1952, p. 3.

- Explanatory arguments which ... account for a phenomenon by reference to statistical laws are not of the strictly deductive type.
- An account of this type will be called an ... inductive explanation.
- C. Hempel, "Aspects of Scientific Explanation", 1965, p. 302.

#### Deductive Inference

#### **Truth Preserving**

- In each possible world:
  - if the premises are true,
  - then the conclusion is true.

#### **Monotonic**

Conclusions are stable in light of further premises.

inference inductive truth preserving, Everything else



deductive

monotonic.

inference

#### deductive

- Calculation
- Refuting universal H
- Verifying existential H
- Deciding between universal H, H'
- Predicting E from H
- Hypotheses compatible with E



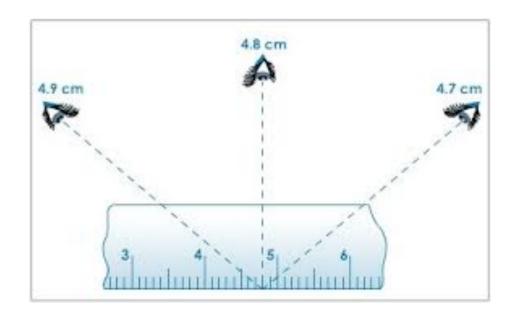
#### inductive

- Inferring universal H
- Choosing between universal H<sub>0</sub>, H<sub>1</sub>, H<sub>2</sub>, ...



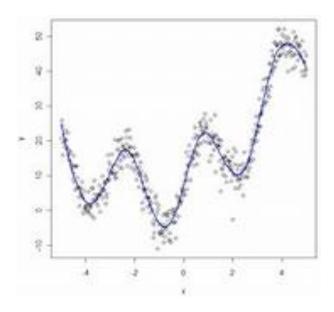
### Real Data

- All real measurements are subject to probable error.
  - It can be reduced through redundancy (sample size).



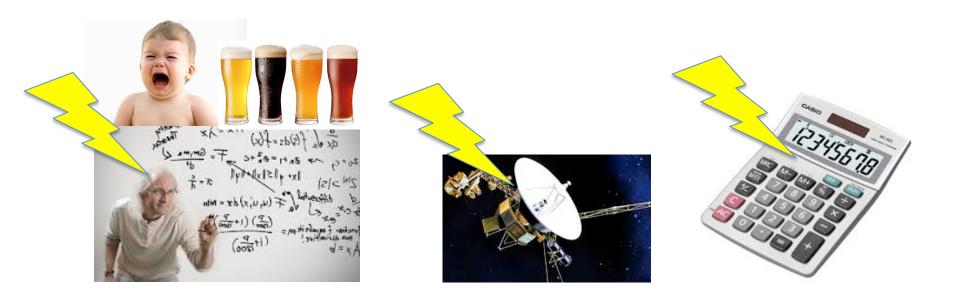
#### Real Predictions

- The predictions of probabilistic theories are subject to probable error.
  - It can be reduced through repeated sampling.



### Real Calculations

- All real calculations are subject to probable error.
  - It can be reduced by redundant codes, circuits, and refereeing.



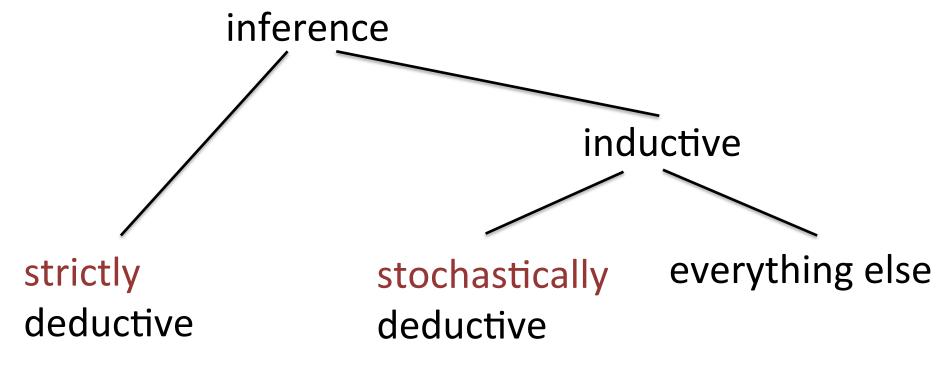
#### Stochastic Deductive Inference

#### Truth preserving in chance

- In each possible world:
  - if the premises are true,
  - then the chance of drawing an erroneous conclusion is low.

#### Monotonic in chance

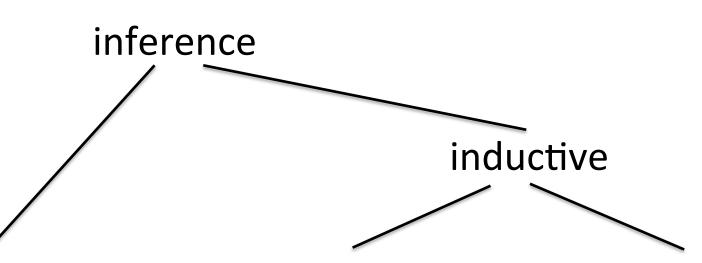
 The chance of producing a conclusion is guaranteed not to drop by much.











#### strictly

#### deductive

- 1. Ideal calculation -----
- 2. Refuting universal  $H_0$  -----
- 3. Verifying existential  $H_1$
- 4. Deciding between universal ---  $H_0$ ,  $H_1$
- 5. Predicting E from H
- 6. Hypotheses compatible with *E*

#### stochastically

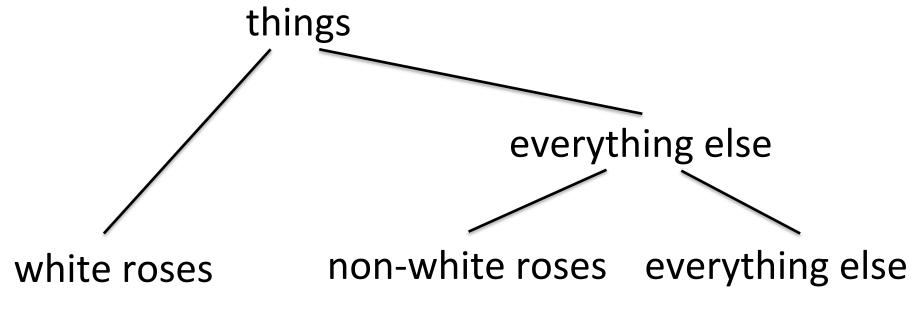
#### deductive

- 1. Real calculation
- 2. Refuting point null  $H_0$
- 3. Verifying composite  $H_1$
- 4. Deciding between point hypotheses  $H_0$ ,  $H_1$
- 5. Direct inference of E from H
- 6. Non-rejection.

#### everything else

- 1. Inferring universal  $H_0$
- 2. Choosing between universal  $H_0$ ,  $H_1$ ,  $H_1$ , ...
- 1. Inferring simple  $H_0$
- Model selection

# **Bad Taxonomy**

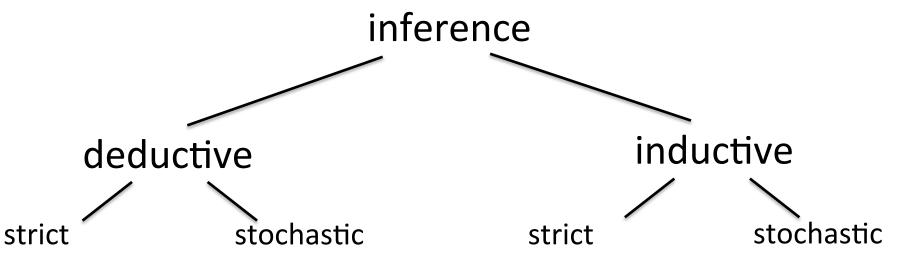








# Improved Taxonomy of Inference



- 1. Refuting universal  $H_{
  m 0}$
- 2. Verifying existential  $H_1$
- 3. Deciding between universal  $H_0$ ,  $H_1$
- 4. Predicting E from H
- 5. Compatibility with *E*
- 6. Ideal calculation

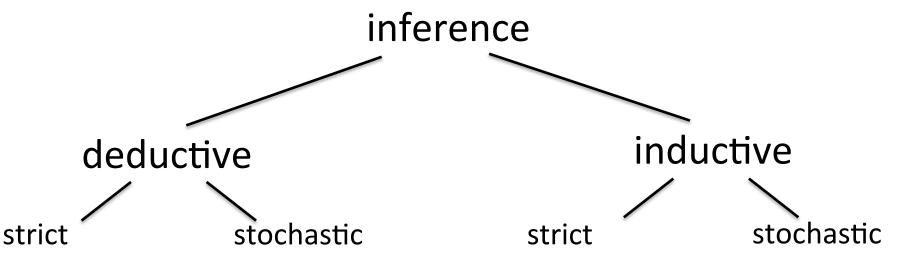
- 1. Refuting point null  $H_0$
- 2. Verifying composite  $H_1$
- 3. Deciding between point hypotheses  $H_0$ ,  $H_1$
- 4. Direct inference of E from H
- 5. Non-rejection.
- 6. Real calculation



- Inferring universal  $H_0$  1. Infe
- 2. Choosing between universal  $H_0$  ,  $H_1$  , ...
- 1. Inferring simple  $H_0$
- 2. Model selection



# Improved Taxonomy of Inference



- Refuting universal  $H_0$
- Verifying existential  $H_1$
- Deciding between universal  $H_0$ ,  $H_1$
- Predicting E from H
- Compatibility with E
- Ideal calculation

- Refuting point null  $H_0$
- 2. Verifying composite  $H_1$
- Deciding between point hypotheses  $H_0$ ,  $H_1$
- Direct inference of *E* from *H*
- Non-rejection. 5.
- Real calculation



- Inferring universal  $H_0$ Inferring simple  $H_0$
- Choosing between universal  $H_0$ ,  $H_1$ ,  $H_1$ , ...
- Model selection

# Question

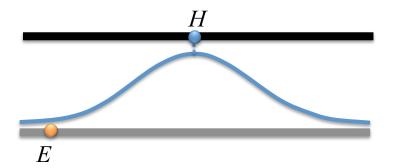
• In strict deduction, the evidence rules out possibilities.



### Question

- In strict deduction, the evidence rules out possibilities.
- In statistical deduction, the sample is logically compatible with every possibility.

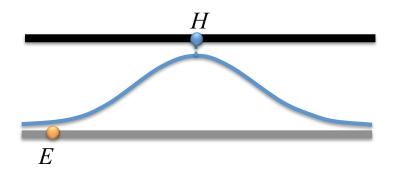




### Question

- In strict deduction, the evidence rules out possibilities.
- In statistical deduction, the sample is logically compatible with every possibility.
- Is there a common, underlying sense of empirical information?





## The Topology of Information

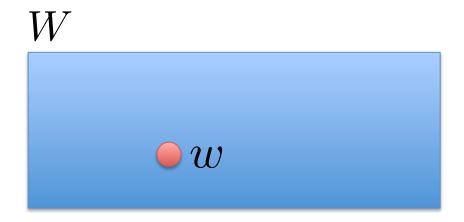


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### Worlds

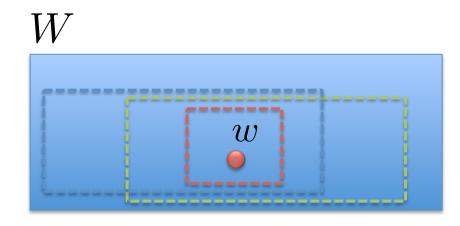
• The points in W are possible worlds.



### The Structure of Information

An **information basis** I is a countable set of information states such that in every world:

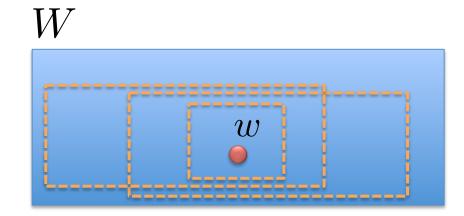
- 1. some information state true;
- each true pair of information states is entailed by a true information state.



### The Structure of Information

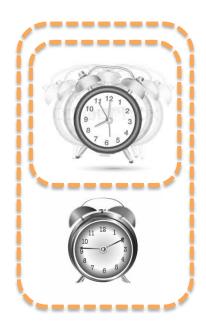
Local information basis at w:

$$\mathcal{I}(w) := \{ E \in \mathcal{I} : w \in E \}.$$



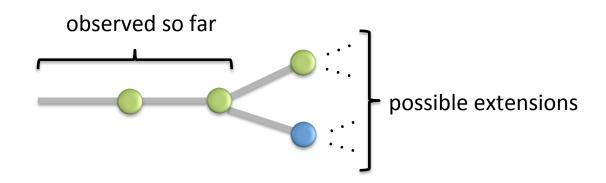
# **Sleeping Beauty Theorist**

 The theorist is awakened by her graduate students only when her theory is refuted.



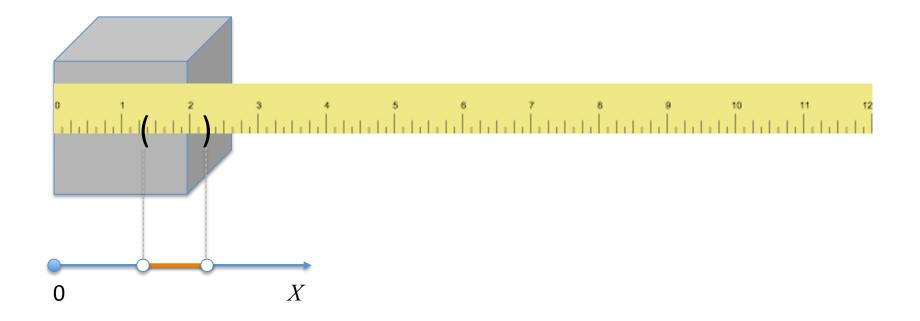
#### Example: Sequential Binary Experiment

**Worlds** = infinite discrete sequences of outcomes. **Information states** = cones of possible extensions:



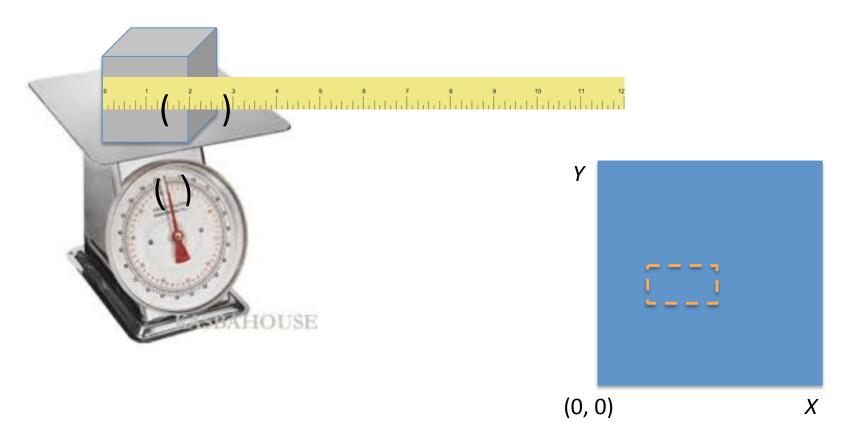
# Example: Measurement of X

- Worlds = real numbers.
- Information states = open intervals.



# Example: Joint Measurement

- Worlds = points in real plane.
- Information states = open rectangles.



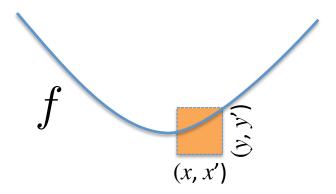
# Example: Equations

• Worlds = functions  $f: \mathbb{R} \to \mathbb{R}$ .



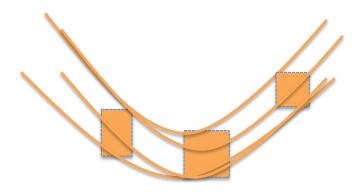
# Example: Laws

• An observation is a joint measurement.



# Example: Laws

 The information state is the set of all worlds that touch each observation.

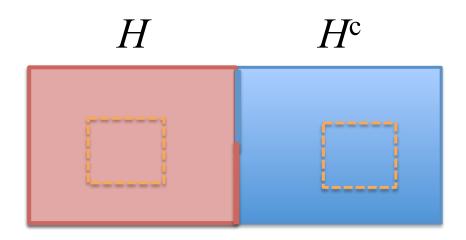


#### Deductive Verification and Refutation

*H* is **verified** by *E* iff  $E \subseteq H$ .

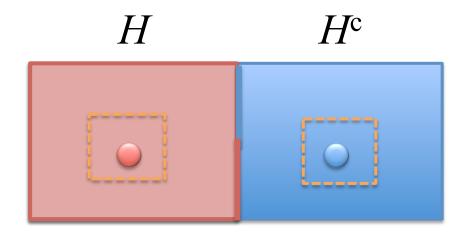
H is **refuted** by E iff  $E \subseteq H^c$ .

H is **decided** by E iff H is either verified or refuted by E.



#### Will be Verified

w is an **interior [exterior] point** of H iff iff H will be verified [refuted] in w iff there is  $E \subseteq \mathcal{I}(w)$  s.t. H is verified [refuted] by E.

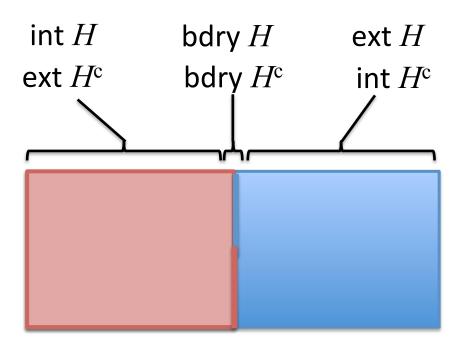


#### Will be Verified

int H := the proposition that H will be verified.

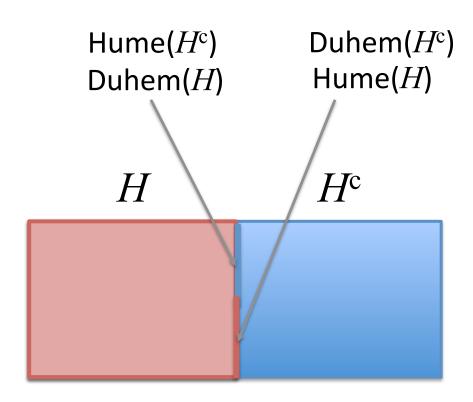
**ext** H := the proposition that H will be refuted.

**bdry** H := the proposition that H will never be decided.



#### Will be Verified

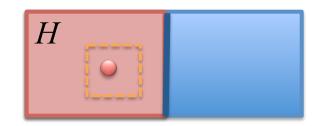
- $bdry(H) \cap H = "you face$ **Hume's problem**w.r.t. <math>H";
- bdry(H)  $\cap H^c$  = "you face **Duhem's problem** w.r.t. H"



## Verifiability, Refutability, Decidability

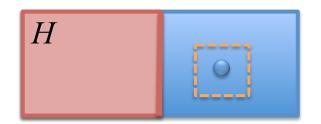
*H* is **verifiable** iff  $H \subseteq int(H)$ .

i.e., iff H will be verified however H is true.

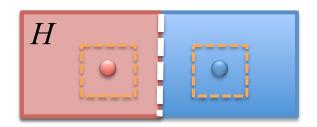


*H* is **refutable** iff  $cl(H) \subseteq H$ .

i.e., iff H will be refuted however H is false.



H is **decidable** iff H is both verifiable and refutable.



- A verification method for H is an inference rule V(E) = A such that in every world w:
  - 1.  $w \in H$ : V converges to H without error.
  - 2.  $w \in H^c$ : V always concludes W.

- A verification method for H is an inference rule V(E) = A such that in every world w:
  - 1.  $w \in H$ : V converges to H without error.
  - 2.  $w \in H^c$ : V always concludes W.
- A **refutation method** for H is just a verification method for  $H^c$ .
- A decision method for H converges to H or to H<sup>c</sup> without error.

- A limiting verification method for H is an inference rule V(E) = A such that in every world w:
  - $w \in H$  iff V converges to some true H' that entails H.
- A **limiting refutation method** for H is a limiting verification method for  $H^c$ .
- A **limiting decision method** for H is a limiting verification method and a limiting refutation for H.

- A verification method for H is an inference rule V(E) = A such that in every world w:
  - 1.  $w \in H$ : V converges to H without error.
  - 2.  $w \in H^c$ : Valways concludes W.
- A **refutation method** for H is just a verification method for  $H^c$ .
- A decision method for H converges to H or to  $H^c$  without error.
- H is methodologically verifiable [refutable, decidable, etc.] iff H has a method of the corresponding kind.

# Verification, Refutation, and Decision are Deductive

**Proposition** (truth preservation).

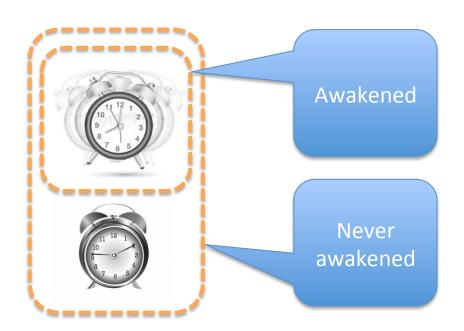
If V is a verifier, refuter or decider for H and V(E) = A, then  $E \subseteq A$ .

#### **Proposition** (monotonicity).

If there is a verifier, refuter or decider for H, then there is a monotonic one that never drops H or  $H^c$  after having concluded it.

# Limiting Verification, Refutation, and Decision are Inductive

**Proposition.** No limiting verifier of "never awakened" is truth preserving or monotonic.



## Topology

Let  $I^*$  denote the closure of I under union.

#### **Proposition:**

If  $(W, \mathcal{I})$  is an information basis then  $(W, \mathcal{I}^*)$  is a topological space.

## Topology

- H is open iff  $H \in \mathcal{I}^*$ .
- H is **closed** iff  $H^c$  is open.
- H is clopen iff H is both closed and open.

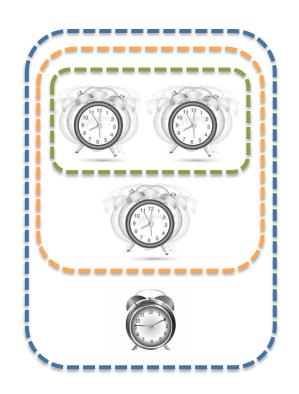
H is locally closed iff H is a difference of open sets.

## Sleeping Theorist Example

 $H_2$  = "Awakened twice" is open.

 $H_1$  = "Awakened once" is locally closed.

 $H_0$  = "Never awakened" is closed.

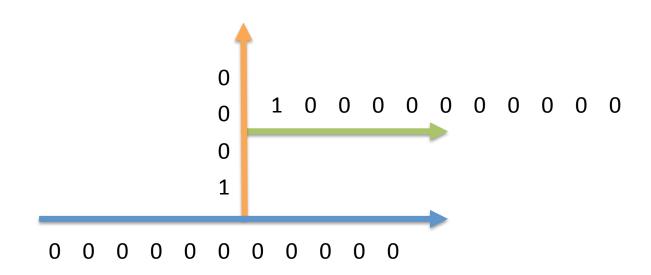


## Sequential Example

```
H_2 = "You will see 1 exactly twice" is open.
```

 $H_1$  = "You will see 1 exactly once" is locally closed.

 $H_0$  = "You will never see 1" is closed.



## **Equation Example**

```
H_2 = "quadratic" is open.

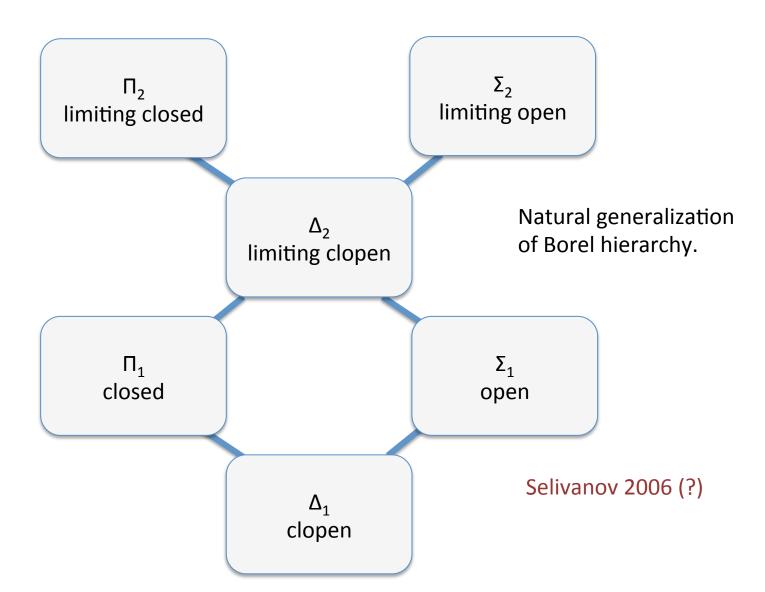
H_1 = "linear" is locally closed.

H_0 = "constant" is closed.
```

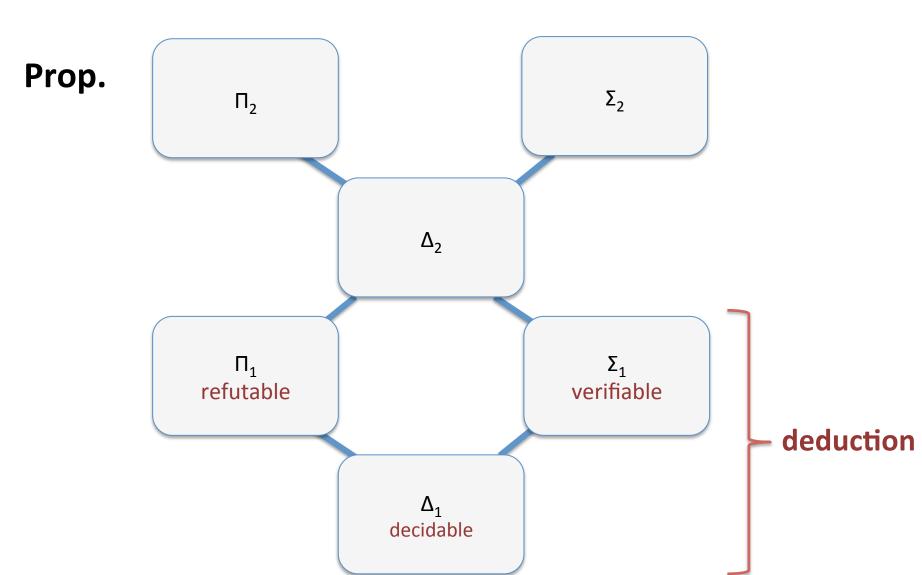
## Topology

- H is limiting open iff H is a countable union of locally closed sets.
- H is **limiting closed** iff  $H^c$  is limiting open.
- H is limiting clopen iff H is both limiting open and limiting closed.

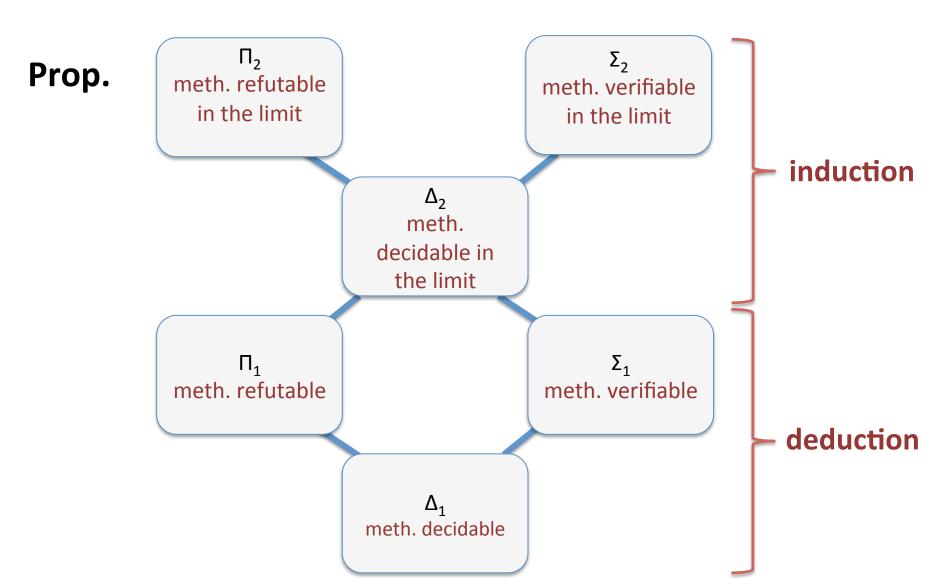
# deBrecht Hierarchy



# **Topology and Pragmatics**



## Topology and Methodology





## **OCKHAM'S TOPOLOGICAL RAZOR**

## Popper Was Doing Topology!

Popper's simplicity relation:

$$A \leq B \Leftrightarrow A \subseteq \mathsf{cl}B.$$

$$H_1 \leq H_2 \leq H_3$$
.

## A Slight Revision

#### Our simplicity relation:

$$A \triangleleft B \Leftrightarrow A \cap \mathsf{cl}(B) \setminus B \neq \varnothing.$$

$$H_1 \triangleleft H_2 \triangleleft H_3$$
.

#### Ockham's Razor

- A question partitions W into possible answers.
- A relevant response is a disjunction of answers.

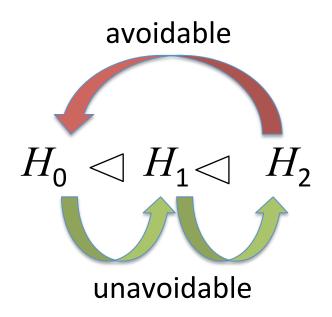
**Proposition.** The following principles are **equivalent**.

- 1. Infer a simplest relevant response in light of E.
- 2. Infer a **refutable** relevant response compatible with E.
- 3. Infer a relevant response that is **not more complex** than the true answer.

## **Epistemic Mandate for Ockham's Razor**

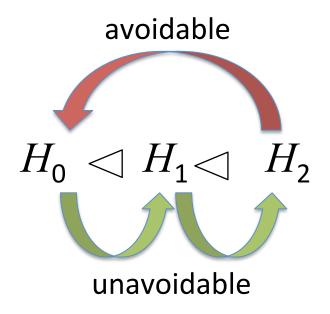
#### If you violate Ockham's razor then

- 1. either you fail to converge to the truth or
- 2. nature can force you into an avoidable cycle of opinions.



## Does Not Presuppose Simplicity

Indeed, by **favoring** a **complex** hypothesis, you incur the avoidable cycle in a **complex** world!

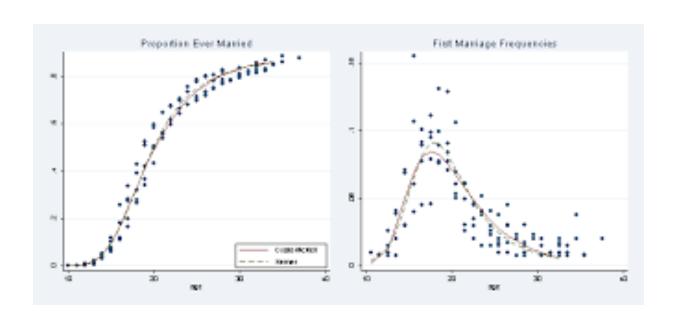




# STATISTICAL INFORMATION TOPOLOGY

## **Statistical Information Topology**

= the topology that lifts the preceding results to statistical inference.



## Skepticism

The above account...

"may be okay if the candidate theories are **deductively** related to observations, but when the relationship is **probabilistic**, I am **skeptical** ...".

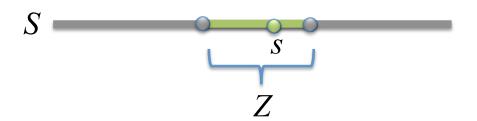


Eliott Sober, Ockham's Razors, 2015

## Epistemology of the Sample

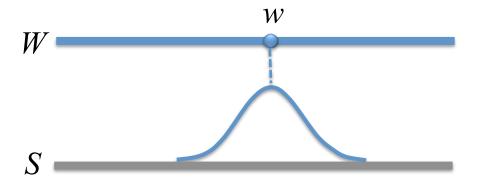
- The sample space S always comes with its own topology  $\mathcal{T}$ .
- $\mathcal{T}$  reflects what is verifiable about the **sample** itself.

s definitely falls within open interval Z.



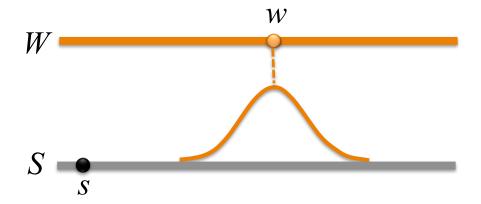
### **Statistics**

• Worlds are probability measures over  $\mathcal{T}$ .



## The Difficulty

- Every sample is logically consistent with all worlds!
- So it seems that statistical information states are all trivial!



## Response

Solve for the unique topology such that:

statistically verifiable = open.

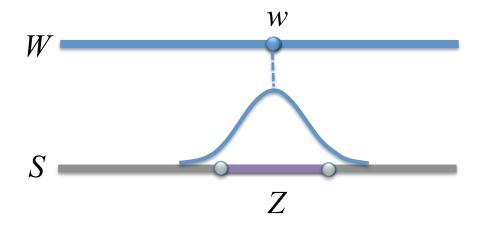
Topology



**Statistics** 

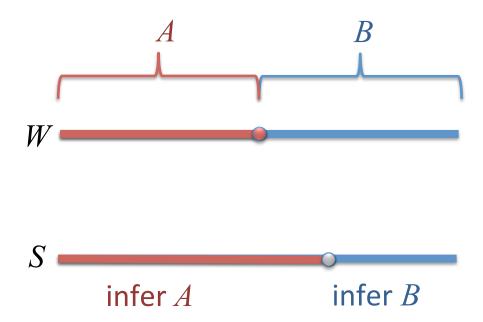
## Feasible Sample Events

- It's impossible to tell whether a point right on the boundary of Z is in or out of Z.
- Z is feasible iff the chance of its boundary is zero in every world.



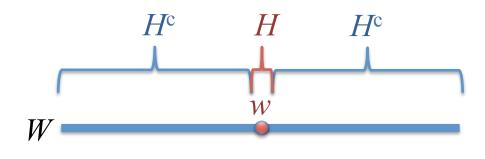
### Feasible Method

A feasible method M is a measurable function from samples to propositions over W such that  $M^{-1}(A)$  is feasible, for all A.



#### Feasible Tests

A feasible test of H is a feasible method that outputs  $H^c$  or W.

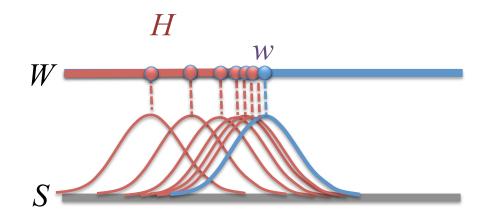




## Statistical Information Topology

 $w \in \operatorname{cl} H$  iff there exists sequence  $(w_n)$  in H, such that for all feasible tests M:

$$\lim_{n\to\infty} p_{w_n}(M \text{ rejects}) \to p_w(M \text{ rejects}).$$



## Weak Topology

**Proposition:** If  $\mathcal{T}$  has a countable basis of feasible regions, then:

statistical information topology = weak topology.

## Weak Topology

**Proposition:** If  $\mathcal{T}$  is second-countable and metrizable, then the weak topology is second-countable and metrizable e.g., by the Prokhorov metric.

• A statistical verification method for H at level  $\alpha > 0$  is a sequence  $(M_n)$  of feasible tests of  $H^c$  such that for every world w and sample size n:

- 1. if  $w \in H$ :  $M_n$  converges in probability to H;
- 2. If  $w \in H^c$ :  $M_n$  concludes W with probability at least 1- $\alpha$ .

• H is **statistically verifiable** iff H has a statistical verification method at each  $\alpha > 0$ .

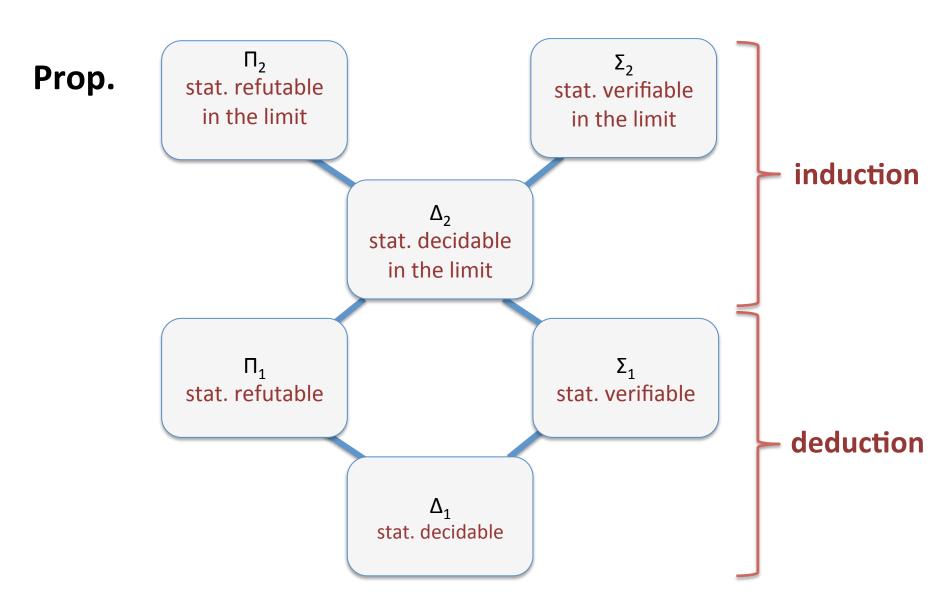
- A statistical verification method for H at level  $\alpha > 0$  is a sequence  $(M_n)$  of feasible tests of  $H^c$  such that for every world w and sample size n:
  - 1. if  $w \in H$ :  $M_n$  converges in probability to H;
  - 2. If  $w \in H^c$ :  $M_n$  concludes W with probability at least  $1-\alpha_n$ ,

for  $\alpha_n \rightarrow 0$ , and dominated by  $\alpha$ .

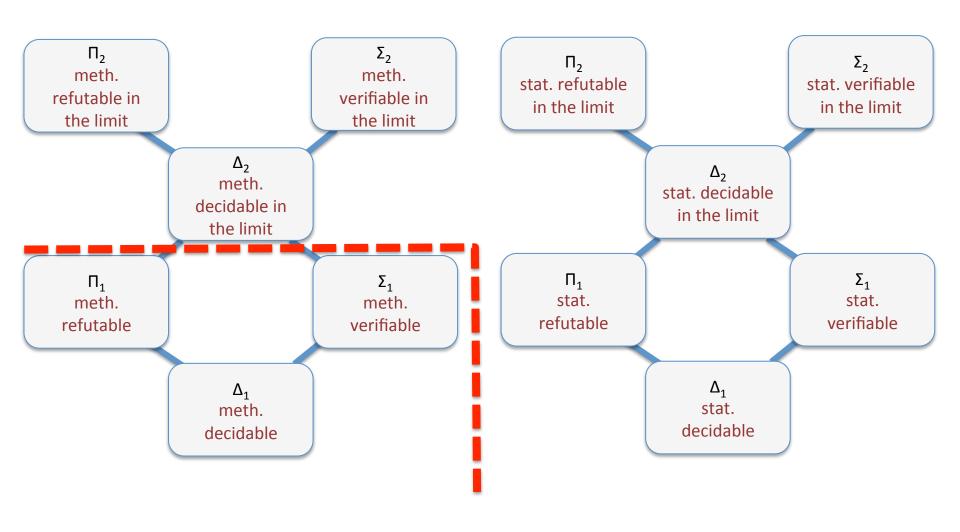
#### Methods

- A limiting statistical verification method for H is a sequence  $(M_n)$  of feasible methods such that:
  - $w \in H$  iff M converges in probability to a true H' that entails H.
- A **limiting statistical refutation method** for H is a limiting verification method for  $H^c$ .
- A **limiting statistical decision method** for H is a limiting verification method and a limiting refutation for H.

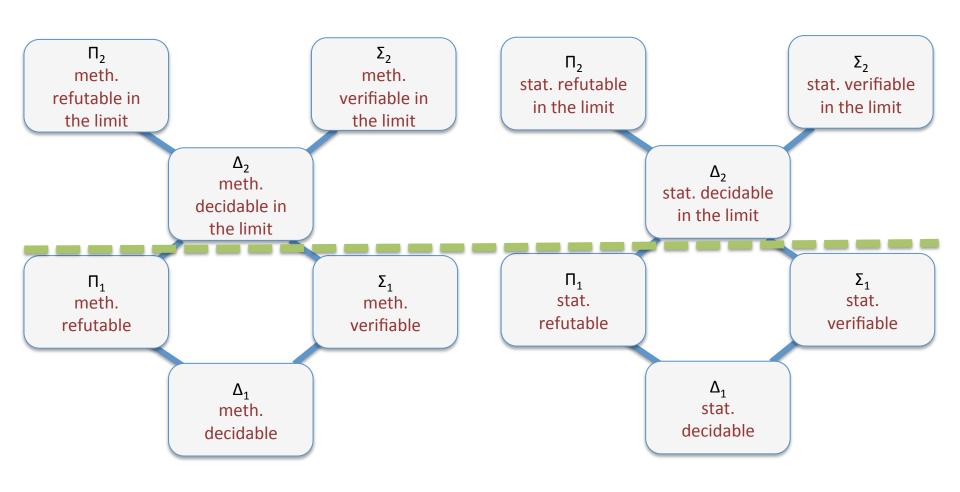
#### Topology and Statistical Methodology



## Deduction vs. Induction: Wrong



## Deduction vs. Induction: Right



## Monotonicity

**Conjecture:** For any open H and  $\alpha > 0$ , there exists a verification method at level  $\alpha$  such that if  $w \in H$ :

$$p_w^{n_2}(M_{n_2} = H) - p_w^{n_1}(M_{n_1} = H) < \alpha,$$

for  $n_2 > n_1$ .



## **Topological Simplicity**

It still makes sense in terms of statistical information topology!

$$A \lhd B \Leftrightarrow A \cap \mathsf{cl}(B) \setminus B \neq \varnothing.$$

$$H_1 \triangleleft H_2 \triangleleft H_3$$
.

#### Ockham's Statistical Razor

Concern: "compatibility with E" is no longer meaningful.

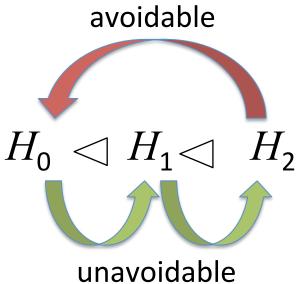
**Response:** the third formulation of O.R. does not mention compatibility with experience!

3. Infer a relevant response that is more complex than the true answer with chance  $< \alpha$ .

#### **Epistemic Mandate for Ockham's Razor**

If you violate Ockham's razor with chance  $\alpha$ , then

- 1. either you fail to converge to the truth in chance or
- 2. nature can force you into an  $\alpha$ -cycle of opinions (complex-simple-complex), even though such cycles are avoidable.



### A New Objective Bayesianism

How much prior bias toward simple models is necessary to avoid  $\alpha$ -cycles?

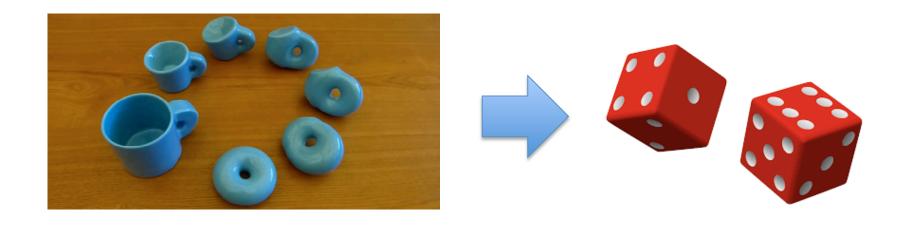
Indifference = ignorance.

truth-conduciveness.

#### **CONCLUSION**

### A Method for Methodology

- 1. Develop basic methodological ideas in topology.
- 2. Port them to statistics via statistical information topology.



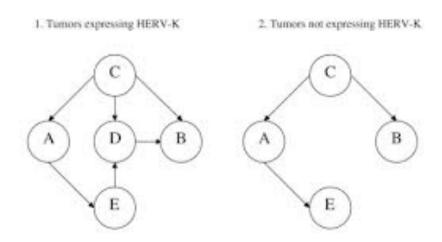
### Some Concluding Remarks

- 1. Information topology is the structure of the scientist's problem context.
- 2. The apparent **analogy** between statistical and ideal methodology reflects **shared topological structure**.
- 3. Thereby, ideal logical/topological ideas can be ported in a direct and uniform fashion to statistics.
- 4. The result is a new, systematic, **frequentist** foundation for **inductive inference** and **Ockham's razor**.

#### ETC.

# Application: Causal Inference from Non-experimental Data

- Causal network inference from retrospective data.
- That is an inductive problem.
- The search is strongly guided by Ockham's razor.
- We have the only non-Bayesian foundation for it.



#### **Application: Science**

- All scientific conclusions are supposed to be counterfactual.
- Scientific inference is strongly simplicity biased.
- Standard ML accounts of Ockham's razor do not apply to such inferences (J. Pearl).
- Our account does.